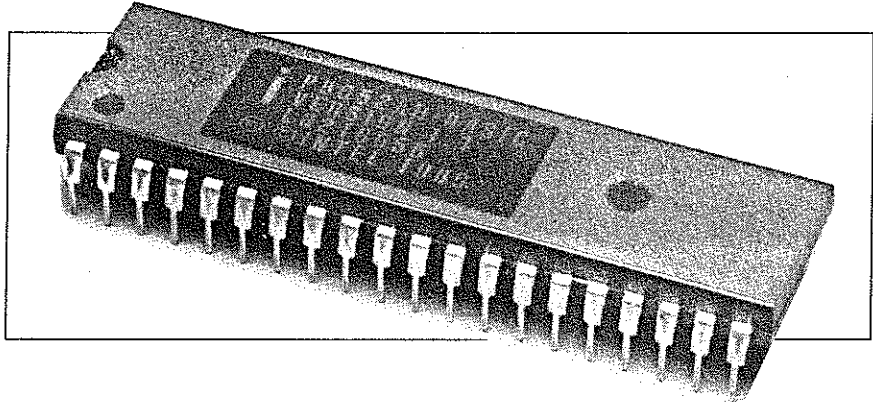


UPGRADE FOR MCS BASIC-52 V.1.1 (Part 2)

Following last month's dealings with the floating-point nucleus and the hex-to-BCD conversion routine in the MCS-52 BASIC interpreter, the authors now tackle some problems with multiplications.



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Continued from the October 1991 issue.

INCONSISTENCIES in the multiplication of two numbers as performed by version 1.1 of Intel's MCS BASIC-52 interpreter can be demonstrated by running the following three small programs, all of which produce wrong results.

```
10 a=1.E-65
20 b=1.E-65
30 ?a*b          (result: 1.0E+126)
```

```
10 a=1.E-65
20 b=1.E-64
30 ?a*b          (result: 0)
```

```
10 a=1.E-64
20 b=1.E-64
30 ?a*b          (result: 1.0E-0)
```

In all three cases, the interpreter should have produced

ERROR: ARITH. UNDERFLOW - IN LINE 30

The above examples point to inconsistencies in limit cases when two numbers are multiplied. On investigating the operation of Intel's BASIC interpreter, a multiplication algorithm of the type shown in Fig. 1 was found. Apparently, the inconsistencies brought to our attention by the above example programs were caused by the exponent adjustment routine, which is listed in Fig. 2. Unfortunately, the errors found in this routine cannot be corrected in BASIC-52 machine code, because some expansion of the machine code is in order. This means that a number of lines should be added to the assembler source file for compensation.

A problem occurs when the result of the instruction SUBB A,#81H is 0FFH, which

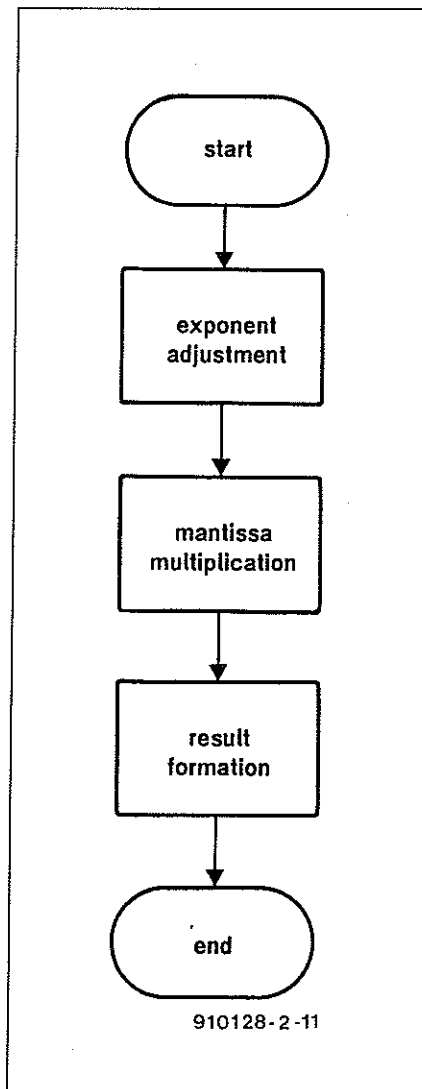


Fig. 1. Flowchart of the routine that handles multiplication of two numbers.

equals exponent E+127. If the program part for mantissa multiplication includes a result that begins with a 0 after the decimal point, the exponent is not incremented but remains

```

1A9AH 9175 ACALL 1C75H
1A9CH 8E0002 CJNE R6,#00H,1AA1H
1A9FH 61B8 AJMP 1B88H
1AA1H 8D2F MOV 2FH,R5
1AA3H EF MOV A,R7
1AA4H 60F9 JZ 1A9FH
1AA6H 2E ADD A,R6
1AA7H 20E705 JB ACC.7,1AAEH
1AAAH 10D706 JBC CY,1AB3H
1AADH 61B2 AJMP 1B82H
1AAEH 5002 JNC 1AB3H
1AB1H 61A1 AJMP 1BA1H
1AB3H 9481 SUBB A,#81H
1AB5H FE MOV R6,A
1AB6H 7188 ACALL 1B88H
1AB8H 7B04 MOV R3,#04H
1ABAH AC01 MOV R4,#01H
1ABCH 8C01 MOV 01H,R4
1ABEH E3 MOVX A,#01H
1ABFH FA MOV R2,A
1B30H 7834 MOV R0,#34H
1B32H E6 MOV A,#00H
1B33H FE MOV R6,A
1B34H 6003 JZ 1B39H
1B36H 717F ACALL 1B7FH
1B38H 18 DEC R0
1B39H 08 INC R0
1B3AH 7408 MOV A,#08H
1B3CH F9 MOV R1,A
1B3DH 28 ADD A,R0
1B3EH F8 MOV R0,A
1B3FH 860500 CJNE 0R0,#05H,1B42H
1B42H 4013 JC 1B57H
1B44H D3 SETB C
1B45H E4 CLR A
1B46H 18 DEC R0
1B47H 36 ADDC A,#00H
1B48H D4 DA A
1B49H D6 XCHD A,#00H
1B4AH 30E409 JNB ACC.4,1B56H
1B4DH 09F5 DJNZ R1,1B44H
1B4FH 18 DEC R0
1B50H 7601 MOV 0R0,#01H
1B52H 717F ACALL 1B7FH
1B54H 8006 SJMP 1B5CH
1B56H 19 DEC R1
1B57H E9 MOV A,R1
1B58H C3 CLR C
1B59H C8 XCH A,R0
1B5AH 98 SUBB A,R0
1B5BH F8 MOV R0,A
1B5CH 792B MOV R1,#2BH
1B5EH E6 MOV A,#00H
1B5FH C4 SWAP A
1B60H 08 INC R0
1B61H D6 XCHD A,#00H
1B62H 4206 CRL 05H,A
1B64H F7 MOV 0R1,A
1B65H 08 INC R0
1B66H 09 INC R1
1B67H B92FF4 CJNE R1,#2FH,1B5EH
1B6AH EE MOV A,R5
1B6BH 7003 JNZ 1B70H
1B6DH 753000 MOV 30H,#00H
1B7FH 0530 INC 30H
1B91H E530 MOV A,30H
1B83H 70F9 JNZ 1B7EH
1B85H D0E0 POP ACC
1B87H D0E0 POP ACC
1B89H 61A1 AJMP 1BA1H
  
```

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Fig. 2. Original program part for exponent adjustment.

Fig. 3. that corr

```

ORG      1A9AH

; TOS - Top Of arithmetic Stack
; NXTOS - NeXt TOS (position behind)

; TOS_MUL1
; routine for multiplication of two FP numbers, one of which
; is located on TOS and the other on NXTOS (NXTOS*TOS)

TOS_MUL1:ACALL  PREP_MUL; clearing of FP working space and reg. prep.
           CJNE  R6,#0,NXTNM_0 ; is NXTOS equal 0?
TOSM_0:  AJMP  ZERO_MANTISSA ; in case TOS or NXTOS are 0
NXTNM_0: MOV  SIGN,R5 ; result mark in SIGN
           MOV  A,R7 ; is TOS equal 0?
           JZ   TOSM_0
           ADD  A,R6 ; addition of exponent degree (exp. multipl.)
           JB  ACC.7,CMPM_EXP ; exp. with different mark or overflow?
           JBC CY,CORM_EXP ; result is bigger than 0.1
           AJMP UNDERFLOW ; in this case exp. sum is < -127
CMPM_EXP: JNC  CORM_EXP; in this case result is smaller than 0.1
OVER_MD:  AJMP OVERFLOW; in this case exp. sum is > 127
CORM_EXP: CLR  MUL_LIMIT_CASE ; flag of multiplication limit case
           SUBB A,#82H ; exp. multipl. results are within the limits
           INC  A ; deduction with 82H because of reduced
           JNC NMARK_L ; conditions
           SETB MUL_LIMIT_CASE ; limit case
NMARK_L:  MOV  R6,A ; exp. adjustment to real value
           ACALL BCD02.1 ; disassembl. of TOS mantissa in LEN_MANTISSA
           MOV  R3,#LEN_BYTE ; acc. (one number in each byte)
           MOV  R4,R1 ; R1 is pointer of NXTOS
MUL_NOOV: MOV  AR1,R4 ; @R1 contains few numbers and disassembled
           MOVX @R1 ; mantissa is successively multiplied by each
           MOV  R2,A ; of them in argument accumulator (R2 is
           ; auxilliary register)

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; TRANSFER
; routine for adjustment of multiplication and division
; results of FP numbers in argument stack and mantissa
; conversion in BCD packed format in argument accumulator

TRANSFER: MOV  R0,#ASCII_DEC ; pointer to MSB result number
           MOV  A,@R0 ; MSB result number is in acc.
           MOV  R6,A
           JZ   MSB_EQ_0; is MSB result number equal 0?
           ACALL INC_EXP ; exp=exp+1 because from the start it was
           DEC  R0 ; taken that MSB result number is equal 0
MSB_EQ_0: INC  R0 ; pointer to MSB-1 number
           MOV  A,#LEN_MANTISSA ; positioning to LSB-1 number
           MOV  R1,A ; R1 is counter
           ADD  A,R0
           MOV  R0,A ; R0 contains the address of LSB-1 number
           CJNE @R0,#05,ROUND1 ; test of remainder to rounding
ROUND1:  JC   NO_ROUND; in case remainder is smaller than 5 there
C_BCD:  SETB C ; is no rounding
           CLR  A ; carry takes one because of rounding
           DEC  R0 ; transfer of carry into higher byte
           ADDC A,@R0 ; addition of 1 to higher byte
           DA  A ; adjustment to BCD format
           XCHD A,@R0 ; result storing in higher byte
           JNB ACC.4,NO_C_BCD ; transfer in higher BCD nibble?
           DJNZ R1,C_BCD; is @R0 the address of MSB number?
           DEC  R0 ; transfer memorizing in the place MSB*1
           MOV  @R0,#01
           ACALL INC_EXP ; transfer is outside decimal point frame
           SJMP BCD1_2
NO_C_BCD: DEC  R1
NO_ROUND: MOV  A,R1 ; address return in R0 to the first bigger BCD
           CLR  C ; number which is not equal 0
           XCH A,R0
           SUBB A,R0
           MOV  R0,A
BCD1_2:  MOV  A,EXPONENT ; test of exceeding in limit case
           JNB MUL_LIMIT_CASE,N_LCASE
           JZ   UNDER_MD; message about underflow
           INC  A
           JZ   UNDER_MD; message about underflow
N_LCASE: MOV  R1,#ARGUMENT_ACC; repacking of BCD number in one byte
XFER_1:  MOV  A,@R0 ; in two BCD numbers in one byte
           SWAP A ; higher number in BCD packed format is in
           INC  R0 ; upper nibble
           XCHD A,@R0 ; BCD packed format is in ACC
           ORL  ARG,A ; the indication if mantissa is 0 is in R6
           MOV  @R1,A ; packing in argument accumulator mantissa
           INC  R0 ; two BCD packed numbers in one byte
           INC  R1
           CJNE R1,#SIGN,XFER_1
           MOV  A,R5
           JNZ RESULT_MAT_OP_TOS ; is mantissa equal 0?
           MOV  EXPONENT,#0

-----

INC_EXP:  INC  EXPONENT; exp. adjustment by 1 (consequence of
           MOV  A,EXPONENT ; transfer after addition)
           JNZ EXP_OK ; test of exponent exceeding
           POP  ACC ; removal from return address stack
           POP  ACC
           JBC MUL_LIMIT_CASE,UNDER_MD ; in limit case it is
           AJMP OVERFLOW; underflow
UNDER_MD: AJMP UNDERFLOW

```

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OFFH. After the result formation (see Fig. 1) and the serial output (display) routine, this gives an exponent E+126.

To avoid inconsistencies, and eliminate errors in limit case multiplications, the program listed in Fig. 3 was developed. The lines marked by vertical bars in particular correct the limit case multiplication errors.

On the basis of the information contained in this and last month's instalment, the authors developed an improved, error-free floating-point (FP) arithmetic for the 8051, on the basis of the FP nucleus extracted from the 8052AH-BASIC. The new FP nucleus uses a modified way of accessing the arithmetic stack, and has faster, shorter code for a number of algorithms. It allows the length of the mantissa of a FP number to range from 2 to 16 digits, while the entire memory map of the FP arithmetic variables is located in the internal memory space of the 8051 microcontroller.

Fig. 3. Source code of the improved multiplication routine. The vertical bars indicate lines that correct the limit case multiplication errors mentioned in this article.

